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Post-quantum blind signcryption scheme from lattice

Key words: Lattice-based cryptosystem; Blind signcryption; Postquantum computing; Learning with error assumption; Small integer solution assumption

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Motivation

- 1. With the development of quantum computing technology, it is imperative to study cryptographic schemes that are resistant to quantum computing.
- 2. As a promising candidate post-quantum cryptosystem, the lattice-based cryptosystem has attracted increasing attention in academic fields.
- 3. Signcryption can simultaneously complete encryption and signature operations in a logical step, and the calculation and communication cost is lower than that of the traditional signature-then-encryption method.

Main idea

- 1. Design a new blind signcryption function which can not only complete the blind signature and encryption functions in a polynomial step but also resist quantum computing attack.
- 2. Prove the security of the scheme in the standard model.
- 3. Verify the efficiency of the proposed scheme through efficiency analysis and parameter size analysis.

Contribution

- 1. A post-quantum blind signcryption scheme from lattice (PQ-LBSCS) is proposed which can simultaneously guarantee blind signature and encryption.
- 2. The new scheme greatly improves the computational efficiency by reducing the length of the ciphertext and public key.

Method

Blind signcryption algorithm

S
$$A : = \mathbf{B}^{(0)} + \sum_{i \in [\lambda]} h_i \mathbf{B}^{(i)}$$

$$h = H_1(\mu, \mathbf{A}_i)$$

$$h_r = H_{\mathbf{A}_s^{(h)}}(h, \mathbf{r}_1)$$

$$\mathbf{y}' = h_r - \mathbf{A}'\mathbf{y}$$

$$\sigma_{or} \leftarrow \text{SampleDG}(\mathbf{T}_s, \mathbf{A}_s^{(0)}, \mathbf{I}, \mathbf{y}', \mathbf{s}_s)$$

$$A : \mathbf{B}^{(0)} + \sum_{i \in [\lambda]} h_i \mathbf{B}^{(i)}$$

$$\mu = (t^{-1}\mathbf{g} + \mathbf{A}_r \mathbf{c}) \mod q$$

$$\sigma = t(\sigma' - \mathbf{c})$$

$$Decompose \ \sigma = (\sigma_1, \sigma_2)$$

$$b^{\mathsf{T}} = 2(\mathbf{s}^{\mathsf{T}} \mathbf{A}_r^{(u)} \mod q) + \mathbf{e}^{\mathsf{T}}$$

$$+ (0, \sigma_1')^{\mathsf{T}} \mod (2q)$$

$$k' = H_3(\sigma_1, \mathbf{b})$$

$$\mathbf{c}' = \varepsilon_{k'}(M|\sigma_2|\mathbf{r}_1|\mathbf{r}_2)$$

$$\mathbf{c} = (\mathbf{u}, \mathbf{b}, \mathbf{c}')$$
Output $\mathbf{c} = (\mathbf{u}, \mathbf{b}, \mathbf{c}')$

Fig. 1 Blind signcryption scenarios

Major results

1. Sizes of some parameters of the new scheme and relevant schemes

Table 2 Size of some parameters

Scheme	Length of ciphertext	Size of params	Number of hash functions	Size of the public key
SZ	$2mk^2n(\log q)^2$	$(d+3)(\log q)^2$	2	$(mk+1)\log q$
YLM	$mkn^2(\log q)^2$	$mnk\log q$	6	$k(k+1)^2 m^2 \lambda (\log q)^2$
YWM	$lmn^2(\log q)^2$	$m^2 n l \log q$	3	$2m^2\log q$
YHW	$2m^2n(\log q)^2$	$km^3\log q$	3	$nm^2\log q$
PQ-LBSCS	$mn^2(\log q)^2$	$2ln^2\log q$	5	$2n^2\log q$

q: a large prime number; m and k: real numbers; n: a security parameter; l: number of columns of matrix F; d: number of vectors B_i in Sun and Zheng (2018)

Major results (Cont'd)

2. Efficiency comparison between the new scheme and relevant schemes

Table 3 Computational efficiency of schemes

Calcana	Computational efficiency			
Scheme -	Signcryption	Unsigncryption		
SZ	$PIS+5D_s+3C_s$	$4D_{\rm s}$ + $6C_{\rm s}$		
YLM	$PIS+5D_s+3C_s$	$7D_{\mathrm{s}}$ + $7C_{\mathrm{s}}$		
YWM	$PIS+5D_s+9C_s$	$6D_{\mathrm{s}}$ + $4C_{\mathrm{s}}$		
YHW	$PIS+4D_{\rm s}+7C_{\rm s}$	$3D_{\rm s}+6C_{\rm s}$		
PQ-LBSCS	$PIS+4D_{\rm s}+2C_{\rm s}$	$2D_{\mathrm{s}}$ + $4C_{\mathrm{s}}$		

 C_s denotes the multiplication operation, D_s denotes the addition operation, and PIS represents the preimage sampling

Major results (Cont'd)

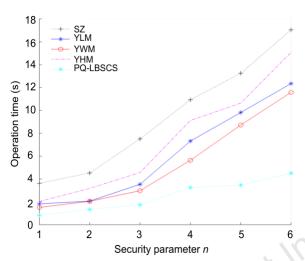


Fig. 2 Signcryption time comparison

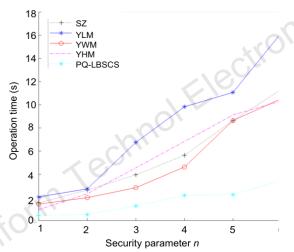


Fig. 3 Unsigncryption time comparison

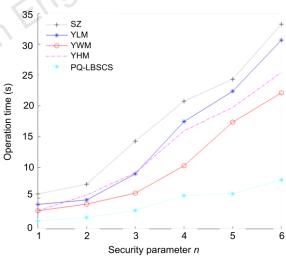


Fig. 4 Time comparison for the whole scheme

Conclusions

- 1. The post-quantum blind signcryption scheme from lattice (PQ-LBSCS) can ensure indistinguishability and unforgeability.
- 2. PQ-LBSCS can blind the message efficiently, achieve anonymous signcryption, and resist quantum computing attackers.
- 3. Simulation results show that the calculation complexity of PQ-LBSCS is relatively low.



Huifang YU was born in Qinghai, China. She received her Ph.D. degree from Shaanxi Normal University. Currently, she is a professor and master supervisor with Xi'an University of Posts & Telecommunications. She has completed more than 10 research projects including a 973 Basic Research Project. She is the PI of more research projects including the National Natural Science Foundation of China. She has published two books and more than 60 papers. She has been authorized five national invention patents. Her main research interests include cryptography and information security.



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