



## An algorithm for identifying symmetric variables in the canonical OR-coincidence algebra system\*

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**Abstract:** To simplify the process for identifying 12 types of symmetric variables in the canonical OR-coincidence (COC) algebra system, we propose a new symmetry detection algorithm based on OR-NXOR expansion. By analyzing the relationships between the coefficient matrices of sub-functions and the order coefficient subset matrices based on OR-NXOR expansion around two arbitrary logical variables, the constraint conditions of the order coefficient subset matrices are revealed for 12 types of symmetric variables. Based on the proposed constraints, the algorithm is realized by judging the order characteristic square value matrices. The proposed method avoids the transformation process from OR-NXOR expansion to AND-OR-NOT expansion, or to AND-XOR expansion, and solves the problem of completeness in the  $d_j$ -map method. The application results show that, compared with traditional methods, the new algorithm is an optimal detection method in terms of applicability of the number of logical variables, detection type, and complexity of the identification process. The algorithm has been implemented in C language and tested on MCNC91 benchmarks. Experimental results show that the proposed algorithm is convenient and efficient.

**Key words:** Symmetric variable,  $d_j$ -map, Canonical OR-coincidence algebra system, Boolean function

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### 1 Introduction

The symmetry of logical variables is a very important property. Symmetric variables exist in most Boolean functions. The symmetry of logical variables has been widely used in the fields of logical synthesis and optimization, ordered binary decision diagram (OBDD) simplification, fault detection, circuit analysis, construction of cryptographic functions, etc.

(Heinrich-Litan and Molitor, 2000; Rice and Muzio, 2002; Rahaman *et al.*, 2003; Acharya *et al.*, 2011; Peng *et al.*, 2011; Shpilka and Tal, 2011; Blais *et al.*, 2012; Kowshik and Kumar, 2013). Detection of the symmetric variables is the basis for all types of applications. There are three main algebra systems of Boolean functions, namely the Boolean logic algebra system based on AND-OR-NOT expansion (Mukhopadhyay, 1963), the canonical Reed-Muller (CRM) algebra system based on AND-XOR expansion (Wu *et al.*, 1982), and the canonical OR-coincidence (COC) algebra system based on OR-NXOR expansion (Cheng *et al.*, 2003). There are also the graphic method (Mukhopadhyay, 1963), the spectral coefficient methods (Hurst, 1977; Rahardja and Falkowski, 1998; Kannurao and Falkowski, 2002), and the tabular method (Lian *et al.*, 2005), which are all used to identify symmetric variables in the traditional AND-OR-NOT algebra system. Since

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Reed (1954) and Muller (1954) proposed the CRM algebra system, the detection algorithm has been researched in the AND-XOR expansion (Falkowski and Kannurao, 1999; Kannurao and Falkowski, 2003). Cheng et al. (2003) proposed the COC algebra system which has the advantages of fault detection and simplified design. Zhao et al. (2006) discussed the detection method based on the  $d_j$ -map in the COC algebra system. There are 12 basic types of symmetry based on a relationship between four co-factors of the Boolean function around two arbitrary variables  $x_i$  and  $x_j$ . The  $d_j$ -map method can be used to identify the type of  $N(x_i|x_j)$  symmetry, but cannot identify the type of  $CN(x_i|x_j)$ ,  $E(x_i|x_j)$ ,  $CE(x_i|x_j)$ ,  $S(x_i|x_j)$ ,  $CS(x_i|x_j)$ ,  $S(x_j|x_i)$ ,  $CS(x_j|x_i)$ ,  $S(x_i|\bar{x}_j)$ ,  $CS(x_i|\bar{x}_j)$ ,  $S(x_j|\bar{x}_i)$ , or  $CS(x_j|\bar{x}_i)$  symmetry. Up to now, the research into symmetric detection has not been perfected for the COC algebra system. To simplify the process for identifying the 12 types of symmetric variables in the COC algebra system, a new symmetry detection algorithm based on OR-NXOR expansion is proposed.

## 2 OR-NXOR expansion of the Boolean function and its coefficient feature

In the COC algebra system, the expansion of the  $n$ -variable Boolean function can be written as

$$f(x_1 \sim x_n) = (d_0 + x_1^1 + x_2^1 + \dots + x_n^1) \odot (d_1 + x_1^0 + x_2^1 + \dots + x_n^1) \odot \dots \odot (d_{2^n-1} + x_1^0 + x_2^0 + \dots + x_n^0), \quad (1)$$

where  $x_i^0 = 0$ ,  $x_i^1 = x_i$ ,  $d_j \in \{0, 1\}$  ( $j=0, 1, \dots, 2^n-1$ ), and ‘ $\odot$ ’ represents the coincidence operation.

Take a three-variable Boolean function as an example:

$$f(x_1 \sim x_n) = (d_0 + x_1 + x_2 + x_3) \odot (d_1 + x_1 + x_2) \odot (d_2 + x_1 + x_3) \odot (d_3 + x_1) \odot (d_4 + x_2 + x_3) \odot (d_5 + x_2) \odot (d_6 + x_3) \odot d_7. \quad (2)$$

To satisfy the correspondence between the subscript of coefficients and the variables in the sum of the terms, expansion (1) can be rewritten as

$$f(x_1 \sim x_n) = (h_{2^n-n} + x_1^1 + x_2^1 + \dots + x_n^1) \odot (h_{2^n-n} + x_1^0 + x_2^1 + \dots + x_n^1) \odot \dots \odot (h_0 + x_1^0 + x_2^0 + \dots + x_n^0). \quad (3)$$

Based on Eqs. (1) and (3), the following transformational relationship between coefficients  $d_j$  and  $h_j$  of a three-variable Boolean function can be obtained: If we transform the subscript of  $d_j$  to the corresponding binary number, the corresponding variables of the 0 value coding can be used to obtain the subscript of  $h_j$ . For example, the binary subscript of  $d_1$  is ‘001’, and the corresponding variables of the 0 value coding are  $x_1$  and  $x_2$ . Thus, the corresponding coefficient  $h_j$  is represented as  $h_{12}$ .

Using Shannon’s law, and considering an arbitrary Boolean function  $f(x_1, \dots, x_i, \dots, x_j, \dots, x_n)$ , the decomposition of the function around two arbitrary variables  $x_i$  and  $x_j$  ( $1 \leq i < j \leq n$ ) is given by

$$f = \bar{x}_i \bar{x}_j f_{00}(x_1, \dots, 0, \dots, 0, \dots, x_n) + \bar{x}_i x_j f_{01}(x_1, \dots, 0, \dots, 1, \dots, x_n) + x_i \bar{x}_j f_{10}(x_1, \dots, 1, \dots, 0, \dots, x_n) + x_i x_j f_{11}(x_1, \dots, 1, \dots, 1, \dots, x_n). \quad (4)$$

**Definition 1** The OR-NXOR expansion coefficient matrices of the sub-functions  $f_{00}(x_1, \dots, 0, \dots, 0, \dots, x_n)$ ,  $f_{01}(x_1, \dots, 0, \dots, 1, \dots, x_n)$ ,  $f_{10}(x_1, \dots, 1, \dots, 0, \dots, x_n)$ , and  $f_{11}(x_1, \dots, 1, \dots, 1, \dots, x_n)$  are represented as  $C_0$ ,  $C_1$ ,  $C_2$ , and  $C_3$ , respectively.

**Definition 2** The order OR-NXOR expansion coefficient matrix of the Boolean function  $f(x_1, \dots, x_i, \dots, x_j, \dots, x_n)$ , which is composed of coefficients  $d_j$  with the subscript in an ascending order, is represented as  $C_I$ . The subscript of  $C_I$  is titled the index of the variables in the sum of the terms. Regard  $C_I$  as a block matrix and decompose it to submatrices corresponding to two arbitrary variables  $x_i$  and  $x_j$ . The order OR-NXOR expansion coefficient subset matrices of  $C_I$  ( $i \in I, j \in I$ ),  $C_I$  ( $i \in I, j \notin I$ ),  $C_I$  ( $i \notin I, j \in I$ ), and  $C_I$  ( $i \notin I, j \notin I$ ) are represented as  $C^0$ ,  $C^1$ ,  $C^2$ , and  $C^3$ , respectively.

Take a three-variable Boolean function as an example. Decomposing the order OR-NXOR expansion coefficient matrix around variables  $x_2$  and  $x_3$ , the order OR-NXOR expansion coefficient subset ma-

trices  $C^0, C^1, C^2$ , and  $C^3$  can be obtained as follows:

$$\begin{aligned} C^0 &= [h_{123}, h_{23}] = [d_0, d_4], C^1 = [h_{12}, h_2] = [d_1, d_5], \\ C^2 &= [h_{13}, h_3] = [d_2, d_6], C^3 = [h_1, h_0] = [d_3, d_7]. \end{aligned}$$

**Theorem 1** The  $n$ -variable Boolean function  $f(x_1 \sim x_n)$  which represents the relationship between the OR-NXOR expansion coefficient matrices of the sub-functions and the order OR-NXOR expansion coefficient subset matrices around two arbitrary logical variables  $x_i$  and  $x_j$  satisfies the following equation:

$$[C_0 \ C_1 \ C_2 \ C_3] = [C^0 \ C^1 \ C^2 \ C^3] \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 \end{bmatrix}. \quad (5)$$

**Proof** According to Kannurao and Falkowski (2003), the relationship between the AND-XOR expansion coefficient matrices of the sub-functions and the order AND-XOR expansion coefficient subset matrices around two arbitrary logical variables  $x_i$  and  $x_j$  satisfies the following equation:

$$[C_0 \ C_1 \ C_2 \ C_3] = [C^0 \ C^1 \ C^2 \ C^3] (T^{\otimes 2})^T, \quad (6)$$

where  $T = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}$ . The transformational relationship between the AND-XOR expansion coefficients and the minterm expansion coefficients can be written as

$$b_n = \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}^{\otimes n} * c_n = T^{\otimes n} * c_n, \quad (7)$$

where  $b_n$  represents a column matrix constituted by coefficients  $b_j$ ,  $c_n$  represents a column matrix constituted by minterm expansion coefficients, ‘\*’ represents the corresponding element of the two matrices multiplication and then adding modulo-2, and  $\otimes n$  represents  $n$  times the Kronecker matrix product.

The transformational relationship between the OR-NXOR expansion coefficients and the maxterm expansion coefficients can be written as

$$\begin{aligned} d_n &= \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}^{\otimes n} * M_n = \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}^{\otimes n} * c_n \\ &= (T^T)^{\otimes n} * c_n, \end{aligned} \quad (8)$$

where  $d_n$  represents a column matrix constituted by coefficients  $d_j$ , and  $M_n$  represents a column matrix constituted by maxterm expansion coefficients.  $M_n$  equals  $c_n$  because they are the values of the function under the same input. By comparing Eq. (7) with Eq. (8), and applying  $T^T$  instead of  $T$  in Eq. (7), the relationship between the OR-NXOR expansion coefficient matrices of the sub-functions and the order OR-NXOR expansion coefficient subset matrices around two arbitrary logical variables  $x_i$  and  $x_j$  can be revealed as follows:

$$\begin{aligned} [C_0 \ C_1 \ C_2 \ C_3] &= [C^0 \ C^1 \ C^2 \ C^3] (T^{\otimes 2})^T \\ &= [C^0 \ C^1 \ C^2 \ C^3] ((T^T)^T)^{\otimes 2} \\ &= [C^0 \ C^1 \ C^2 \ C^3] T^{\otimes 2} \\ &= [C^0 \ C^1 \ C^2 \ C^3] \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix}^{\otimes 2} \\ &= [C^0 \ C^1 \ C^2 \ C^3] \begin{bmatrix} 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 1 \end{bmatrix}. \end{aligned}$$

### 3 Symmetric and antisymmetric conditions of logical variables

According to Falkowski and Kannurao (1999) and Kannurao and Falkowski (2003), there are 12 basic types of symmetry, four of which are two-variable symmetry and eight of which are single variable symmetry. Six types of symmetric definitions and six types of antisymmetric definitions are formulated in Tables 1 and 2, respectively.

**Theorem 2** If an  $n$ -variable Boolean function  $f(x_1, \dots, x_i, \dots, x_j, \dots, x_n)$  exists with symmetric variables, then the order OR-NXOR expansion coefficient subset matrices around two arbitrary logical variables  $x_i$  and  $x_j$  satisfy the conditions listed in Table 3.

**Proof** Take the type of  $E(x_i|x_j)$  symmetry as an example. From the symmetric definition in Table 1, if a Boolean function has the type of  $E(x_i|x_j)$  symmetry, then  $f(x_1, \dots, 0, \dots, 0, \dots, x_n) = f(x_1, \dots, 1, \dots, 1, \dots, x_n)$ , and hence  $C_0 = C_3$ . According to Theorem 1,

$$C_0 = C^0 \odot C^1 \odot C^2 \odot C^3, C_3 = C^3,$$

$$\text{or } C^0 \odot C^1 \odot C^2 \odot C^3 = C^3.$$

**Table 1 The six types of symmetry**

Symbol	Symmetric definition
$E(x_i x_j)$	$f(x_1, \dots, 0, \dots, 0, \dots, x_n) = f(x_1, \dots, 1, \dots, 1, \dots, x_n)$
$N(x_i x_j)$	$f(x_1, \dots, 0, \dots, 1, \dots, x_n) = f(x_1, \dots, 1, \dots, 0, \dots, x_n)$
$S(x_i x_j)$	$f(x_1, \dots, 0, \dots, 1, \dots, x_n) = f(x_1, \dots, 1, \dots, 1, \dots, x_n)$
$S(x_i   \bar{x}_j)$	$f(x_1, \dots, 0, \dots, 0, \dots, x_n) = f(x_1, \dots, 1, \dots, 0, \dots, x_n)$
$S(x_j x_i)$	$f(x_1, \dots, 1, \dots, 0, \dots, x_n) = f(x_1, \dots, 1, \dots, 1, \dots, x_n)$
$S(x_j   \bar{x}_i)$	$f(x_1, \dots, 0, \dots, 0, \dots, x_n) = f(x_1, \dots, 0, \dots, 1, \dots, x_n)$

**Table 2 The six types of antisymmetry**

Symbol	Antisymmetric definition
$CE(x_i x_j)$	$f(x_1, \dots, 0, \dots, 0, \dots, x_n) = \overline{f(x_1, \dots, 1, \dots, 1, \dots, x_n)}$
$CN(x_i x_j)$	$f(x_1, \dots, 0, \dots, 1, \dots, x_n) = \overline{f(x_1, \dots, 1, \dots, 0, \dots, x_n)}$
$CS(x_i x_j)$	$f(x_1, \dots, 0, \dots, 1, \dots, x_n) = \overline{f(x_1, \dots, 1, \dots, 1, \dots, x_n)}$
$CS(x_i   \bar{x}_j)$	$f(x_1, \dots, 0, \dots, 0, \dots, x_n) = \overline{f(x_1, \dots, 1, \dots, 0, \dots, x_n)}$
$CS(x_j x_i)$	$f(x_1, \dots, 1, \dots, 0, \dots, x_n) = \overline{f(x_1, \dots, 1, \dots, 1, \dots, x_n)}$
$CS(x_j   \bar{x}_i)$	$f(x_1, \dots, 0, \dots, 0, \dots, x_n) = \overline{f(x_1, \dots, 0, \dots, 1, \dots, x_n)}$

**Table 3 The order OR-NXOR expansion coefficient subset matrices of symmetric variables**

Type of symmetry	Satisfied conditions of the order OR-NXOR expansion coefficient subset matrices
$E(x_i x_j)$	$C^0 \odot C^1 \odot C^2 = [11 \dots 1]^T$
$N(x_i x_j)$	$C^1 \odot C^2 = [11 \dots 1]^T$
$S(x_i x_j)$	$C^1 = [11 \dots 1]^T$
$S(x_i   \bar{x}_j)$	$C^0 \odot C^1 = [11 \dots 1]^T$
$S(x_j x_i)$	$C^2 = [11 \dots 1]^T$
$S(x_j   \bar{x}_i)$	$C^0 \odot C^2 = [11 \dots 1]^T$
$CE(x_i x_j)$	$C^0 \odot C^1 \odot C^2 = [10 \dots 0]^T$
$CN(x_i x_j)$	$C^1 \odot C^2 = [10 \dots 0]^T$
$CS(x_i x_j)$	$C^1 = [10 \dots 0]^T$
$CS(x_i   \bar{x}_j)$	$C^0 \odot C^1 = [10 \dots 0]^T$
$CS(x_j x_i)$	$C^2 = [10 \dots 0]^T$
$CS(x_j   \bar{x}_i)$	$C^0 \odot C^2 = [10 \dots 0]^T$

So,  $C^0 \odot C^1 \odot C^2 = [1 \ 1 \ \dots \ 1]^T$ .

By the same argument, the other 11 symmetry types in Table 3 can be proved.

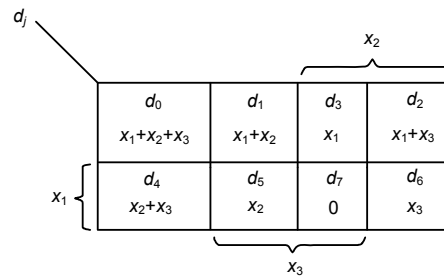
### 4 Algorithm for identifying symmetric variables based on $d_j$ -map

#### 4.1 Principle of the algorithm

**Definition 3** ( $d_j$ -map (Cheng et al., 2003)) The  $d_j$ -map is a representation of a Boolean function in the COC algebra system. Fig. 1 gives the  $d_j$ -map of the three-variable Boolean function.

The  $d_j$ -map has the following features:

1. Each square corresponds to the sum of the terms. If a  $d_j=0$  entry is included within an  $x_i=0$  area, then variable  $x_i$  appears in uncomplemented form in the sum of the terms. On the other hand, if within an  $x_i=1$  area, then variable  $x_i$  does not appear in the sum of the terms.



**Fig. 1 The  $d_j$ -map of a three-variable Boolean function**

2. Adjacent squares differ in one variable.

When the logical variables  $x_i x_j$  equal 00, 01, 10, or 11, the corresponding squares are titled the characteristic squares of the  $d_j$ -map. The filled values of the characteristic squares are titled the characteristic square values.

**Definition 4** (The order characteristic square value matrix  $[d_j]_{x_i x_j}$ ) If the characteristic square values are arranged in ascending order of the subscript of  $d_j$ , then coefficients  $d_j$  constitute the order characteristic square value matrix, represented as  $[d_j]_{x_i x_j}$ . The order characteristic square value matrices are represented as  $[d_j]_{00}$ ,  $[d_j]_{01}$ ,  $[d_j]_{10}$ , and  $[d_j]_{11}$  when  $x_i x_j$  equal 00, 01, 10, and 11, respectively.

From the  $d_j$ -map features and Definitions 1–4, it is easy to obtain the following relationship between the order OR-NXOR expansion coefficient subset

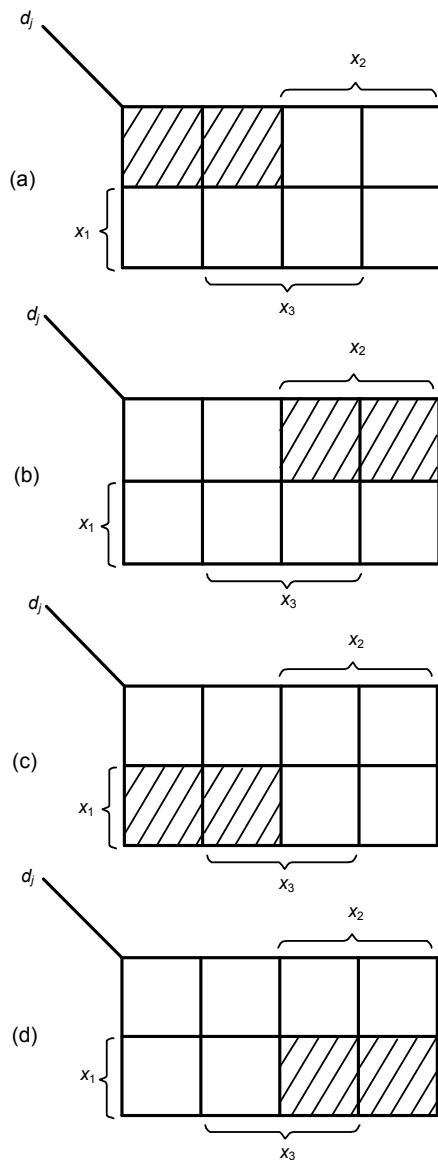
matrices  $C^0, C^1, C^2, C^3$ , and  $[d_j]_{x_i x_j}$ :

$$C^0=[d_j]_{00}, C^1=[d_j]_{01}, C^2=[d_j]_{10}, C^3=[d_j]_{11}.$$

Taking a three-variable Boolean function as an example, the characteristic squares of the  $d_j$ -map around logical variables  $x_1$  and  $x_2$  demonstrated by oblique lines are shown in Fig. 2.

From the above discussions, the following deduction can be obtained:

**Deduction 1** If the order characteristic square value matrix  $[d_j]_{x_i x_j}$  of the  $n$ -variable Boolean function  $f(x_1,$



**Fig. 2** The characteristic squares around logical variables  $x_1$  and  $x_2$

(a)  $x_1 x_2=00$ ; (b)  $x_1 x_2=01$ ; (c)  $x_1 x_2=10$ ; (d)  $x_1 x_2=11$

$\dots, x_i, \dots, x_j, \dots, x_n$ ) satisfies the conditions listed in Table 4, then there exist corresponding symmetric variables.

**Proof** Take the satisfied conditions  $[d_j]_{00} \odot [d_j]_{01} \odot [d_j]_{10}=[1 \ 1 \ \dots \ 1]^T$  as an example. If the order characteristic square value matrix  $[d_j]_{x_i x_j}$  of the  $n$ -variable Boolean function satisfies  $[d_j]_{00} \odot [d_j]_{01} \odot [d_j]_{10}=[1 \ 1 \ \dots \ 1]^T$ , namely

$$C^0 \odot C^1 \odot C^2 = [11 \dots 1]^T, \tag{9}$$

then expansion (9) can be rewritten as

$$C^0 \odot C^1 \odot C^2 \odot C^3 = [11 \dots 1]^T \odot C^3 = C^3.$$

According to Theorem 1,  $C^0 \odot C^1 \odot C^2 \odot C^3=C_0$  and  $C^3=C_3$ . So,  $C_0=C_3$ , namely  $f(x_1, \dots, 0, \dots, 0, \dots, x_n)=f(x_1, \dots, 1, \dots, 1, \dots, x_n)$ . From Table 1, the  $n$ -variable Boolean function exists with the  $E(x_i|x_j)$  symmetry.

By the same argument, the other 11 satisfied conditions of  $[d_j]_{x_i x_j}$  in Table 4 can be proved.

**Table 4** The symmetric conditions of logical variables

The satisfied conditions of $[d_j]_{x_i x_j}$	Existing symmetry type
$[d_j]_{00} \odot [d_j]_{01} \odot [d_j]_{10} = [1 \ 1 \ \dots \ 1]^T$	$E(x_i x_j)$
$[d_j]_{01} \odot [d_j]_{10} = [1 \ 1 \ \dots \ 1]^T$	$N(x_i x_j)$
$[d_j]_{01} = [1 \ 1 \ \dots \ 1]^T$	$S(x_i x_j)$
$[d_j]_{00} \odot [d_j]_{01} = [1 \ 1 \ \dots \ 1]^T$	$S(x_i   \bar{x}_j)$
$[d_j]_{10} = [1 \ 1 \ \dots \ 1]^T$	$S(x_j x_i)$
$[d_j]_{00} \odot [d_j]_{10} = [1 \ 1 \ \dots \ 1]^T$	$S(x_j   \bar{x}_i)$
$[d_j]_{00} \odot [d_j]_{01} \odot [d_j]_{10} = [1 \ 0 \ \dots \ 0]^T$	$CE(x_i x_j)$
$[d_j]_{01} \odot [d_j]_{10} = [1 \ 0 \ \dots \ 0]^T$	$CN(x_i x_j)$
$[d_j]_{01} = [1 \ 0 \ \dots \ 0]^T$	$CS(x_i x_j)$
$[d_j]_{00} \odot [d_j]_{01} = [1 \ 0 \ \dots \ 0]^T$	$CS(x_i   \bar{x}_j)$
$[d_j]_{10} = [1 \ 0 \ \dots \ 0]^T$	$CS(x_j x_i)$
$[d_j]_{00} \odot [d_j]_{10} = [1 \ 0 \ \dots \ 0]^T$	$CS(x_j   \bar{x}_i)$

Based on the above discussions, the method for identifying 12 types of symmetric variables based on  $d_j$ -map for the  $n$ -variable Boolean function is decomposed into the following two steps:

1. Draw the  $d_j$ -map of the  $n$ -variable Boolean function.

2. Identify the 12 types of symmetric variables by judging whether  $[d_j]_{x_i x_j}$  satisfies Deduction 1.

### 4.2 An example of the algorithm

**Example 1** Consider a three-variable Boolean function  $f_1(x_1 \sim x_3) = \ominus \Pi 0, 3, 4$ , and apply the  $d_j$ -map to identify symmetric variables.

The  $d_j$ -map of  $f_1(x_1 \sim x_3)$  is shown in Fig. 3. From the  $d_j$ -map,  $[d_j]_{x_1 x_2}$  is listed as follows:

$[d_j]_{00} = [0 \ 0]^T$ ,  $[d_j]_{01} = [1 \ 0]^T$ ,  $[d_j]_{10} = [0 \ 1]^T$ ,  $[d_j]_{11} = [1 \ 1]^T$ . According to Deduction 1,  $[d_j]_{01} = [1 \ 0]^T$  and then the  $CS(x_1|x_2)$  symmetry exists;  $[d_j]_{00} \odot [d_j]_{10} = [1 \ 0]^T$ , then the  $CS(x_2|\bar{x}_1)$  symmetry exists;  $[d_j]_{00} \odot [d_j]_{01} \odot [d_j]_{10} = [1 \ 1]^T$ , then the  $E(x_1|x_2)$  symmetry exists. By the same argument, the  $CN(x_1|x_3)$ ,  $CS(x_1|\bar{x}_3)$ ,  $S(x_3|\bar{x}_1)$ ,  $CE(x_2|x_3)$ ,  $CS(x_2|\bar{x}_3)$ , and  $S(x_3|x_2)$  symmetries can be identified.

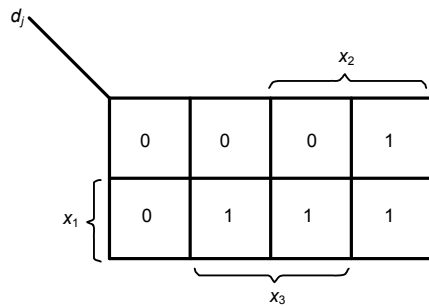


Fig. 3 The  $d_j$ -map of  $f_1$

### 5 Algorithm realization

Based on the discussions on the principle and example of the algorithm, the 12 types of symmetric variables can be identified by judging  $[d_j]_{x_i x_j}$ . The algorithm is decomposed into the following steps:

Step 1: A 2D array  $y$  is defined which is used to conserve the identified function  $f_1(x_1 \sim x_n)$ . Array  $y$  has  $2^n$  rows and  $n+1$  columns. The first  $n$  columns conserve the variables coding of  $f_1(x_1 \sim x_n)$ . The last column conserves coefficients  $d_j$  of  $f_1(x_1 \sim x_n)$ . Taking Example 1 as an example, a 2D array  $y_1$ , which is used to conserve the identified function  $f_1(x_1 \sim x_3)$ , can be represented as shown in Table 5.

Step 2: Define 1D arrays  $C^0$ ,  $C^1$ ,  $C^2$ , and  $C^3$ . Arbitrarily choose two columns of array  $y$ , namely the value coding of  $x_i$  and  $x_j$  ( $1 \leq i < j \leq n$ ), and compare the values of  $y[k][i]$ ,  $y[k][j]$  ( $1 \leq k \leq 2^n$ ) to the characteristic coding.

Table 5 Representation of columns of array  $y_1$

Column 1	Column 2	Column 3	Column 4
0	0	0	0
0	0	1	0
0	1	1	0
0	1	0	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

1. If the values of  $y[k][i]$  and  $y[k][j]$  all equal 0, then store the value of  $y[k][n+1]$  in array  $C^0$ .

2. If the value of  $y[k][i]$  equals 0 and the value of  $y[k][j]$  equals 1, then store the value of  $y[k][n+1]$  in array  $C^1$ .

3. If the value of  $y[k][i]$  equals 1 and the value of  $y[k][j]$  equals 0, then store the value of  $y[k][n+1]$  in array  $C^2$ .

4. If the values of  $y[k][i]$  and  $y[k][j]$  all equal 1, then store the value of  $y[k][n+1]$  in array  $C^3$ .

Step 3: Judge the values of arrays  $C^0$ ,  $C^1$ ,  $C^2$ , and  $C^3$ .

1. The corresponding elements of arrays  $C^0$ ,  $C^1$ , and  $C^2$  carry out the coincidence operation. If all operation results equal 1, then output  $E(x_i|x_j)$ ; If the operation results of the first column equal 1, and the others equal 0, then output  $CE(x_i|x_j)$ .

2. The corresponding elements of arrays  $C^1$  and  $C^2$  carry out the coincidence operation. If all operation results equal 1, then output  $N(x_i|x_j)$ ; If the operation results of the first column equal 1, and the others equal 0, then output  $CN(x_i|x_j)$ .

3. The corresponding elements of  $C^0$  and  $C^1$  carry out the coincidence operation. If all operation results equal 1, then output  $S(x_i|\bar{x}_j)$ ; If the operation results of the first column equal 1, and the others equal 0, then output  $CS(x_i|\bar{x}_j)$ .

4. If all elements of array  $C^2$  equal 1, then output  $S(x_j|x_i)$ ; If the first element of array  $C^2$  equals 1 and the others equal 0, then output  $CS(x_j|x_i)$ .

5. The corresponding elements of  $C^0$  and  $C^2$  carry out the coincidence operation. If all operation results equal 1, then output  $S(x_j|\bar{x}_i)$ ; If the operation results of the first column equal 1, and the others equal 0, then output  $CS(x_j|\bar{x}_i)$ .

6. If all elements of array  $C^1$  equal 1, then output  $S(x_i|x_j)$ ; If the first element of array  $C^1$  equals 1 and the others equal 0, then output  $CS(x_i|x_j)$ .

Step 4: Repeat step 2.

Step 5: Termination the loop.

The flow chart of the program implemented in C language is shown in Fig. 4.

### 6 Experimental results

The proposed algorithm was implemented in C language on the Linux platform and run on an IBM P5 560Q server (eight cores, 64 GB memory, using the Linux enterprise server operating system, SCSI disk). To verify the proposed algorithm, it was tested on MCNC91 benchmarks. The results of the test are shown in Table 6 for single-output functions up to 18 variables. The test function ‘test 1’ is the same as that used in Example 1 in Section 4.2. The results identified by programming the algorithm are consistent with the conclusions identified by judging  $[d_j]_{x_i x_j}$  of the function. The validity of the algorithm is also proven. The symmetry detection time for the same number of logical variables in the test function depends on the number of symmetric variables. The

more the symmetric variables, the longer the detection time required, as indicated by the test functions ‘pde’ and ‘ig9’ in Table 6. The required detection time increases obviously along with the increasing number of logical variables in the test function.

Up to now, there have been no publications on the algorithm which can be programmed for

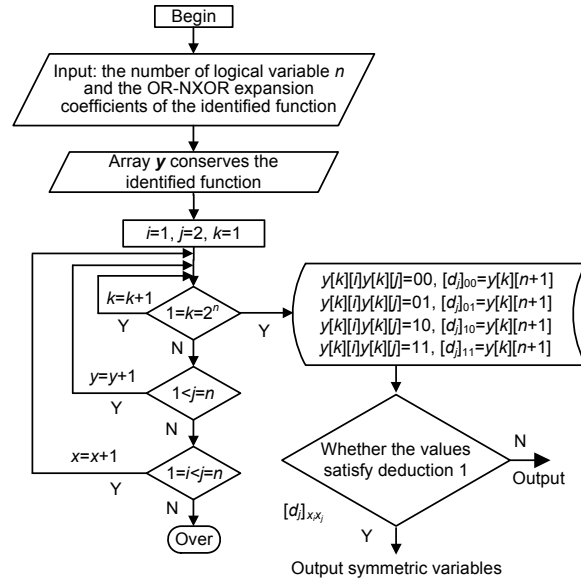


Fig. 4 The flow chart of the program

Table 6 The experimental results for the test function

Name	n	t (s)	Number of existing symmetric variables											
			$E(x_i x_j)$	$N(x_i x_j)$	$S(x_i x_j)$	$S(x_i \bar{x}_j)$	$S(x_j x_i)$	$S(x_j \bar{x}_i)$	$CE(x_i x_j)$	$CN(x_i x_j)$	$CS(x_i x_j)$	$CS(x_i \bar{x}_j)$	$CS(x_j x_i)$	$CS(x_j \bar{x}_i)$
test1	3	1.08	1	0	0	0	1	1	1	1	1	2	0	1
c4	4	3.12	1	0	1	1	0	1	0	1	0	1	0	0
ie3	5	6.28	0	1	1	0	0	0	2	0	0	1	1	0
6fe	6	8.21	1	1	1	1	0	0	0	1	0	1	0	0
7ec	7	9.45	0	0	1	0	1	0	1	0	2	0	1	0
dc8	8	11.22	0	1	0	0	1	0	0	1	0	1	0	1
pde	9	15.26	1	0	1	0	0	1	0	0	0	0	0	1
ig9	9	22.83	1	0	1	0	1	0	1	1	1	0	1	1
qtt	10	24.61	1	1	0	1	1	0	1	0	0	1	0	1
exp	11	27.56	0	2	1	1	1	0	0	1	0	1	0	0
ss1	12	30.8	2	0	0	1	1	2	0	1	0	0	1	1
ort	13	37.83	1	0	1	1	1	0	1	0	1	0	1	0
we2	14	45.2	0	0	1	0	1	0	1	0	0	0	1	0
xxt	15	57.21	1	0	0	0	0	1	0	1	0	0	0	0
uj1	16	72	0	0	0	0	1	1	0	0	0	1	1	0
le5	17	102	0	1	0	0	0	1	0	0	0	1	0	0
vcc	18	148	0	0	0	0	0	0	0	0	0	0	0	0

Name: the test function; n: the number of logical variables in the test function; t: the symmetry detection time (s)

identifying 12 types of symmetric variables in the COC algebra system for comparison.

## 7 Comparison of different identification methods

The performance comparisons of different methods for identifying symmetric variables are listed in Table 7. The graphic method and the  $d_j$ -map method are restricted by the number of logical variables. Generally, these two methods are suitable for the functions where the number of logical variables is no more than six. The process is complicated for calculating spectrum coefficients and separating the spectral coefficients into  $n(n-1)/2$  groups by spectral coefficient methods. The AND-XOR expansion coefficient methods (Falkowski and Kannurao, 1999; Kannurao and Falkowski, 2003) also need to separate the AND-XOR expansion coefficients into  $n(n-1)/2$  groups. It is difficult to classify the coefficients by spectral coefficient methods and the AND-XOR expansion coefficient method when the number of logical variables exceeds six. Only the  $E(x_i|x_j)$  and  $N(x_i|x_j)$  symmetries can be identified by the tabular method. From Table 7, the proposed method is the optimal detection method in terms of the applicability of the number of logical variables, detection type, and the complexity of the identification process. If a Boolean function is expressed as an OR-NXOR expansion, the proposed method which avoids the transformation process from the OR-NXOR expansion to AND-OR-

NOT expansion or to the AND-XOR expansion and solves the completeness of the  $d_j$ -map method, is obviously the most efficient. If a Boolean function is expressed as an AND-OR-NOT expansion or AND-XOR expansion, by the transformation process from the AND-OR-NOT expansion or AND-XOR expansion to the OR-NXOR expansion, the proposed algorithm is also the most appropriate method.

## 8 Conclusions

The Boolean logic algebra system, the CRM algebra system, and the COC algebra system are three algebra systems for Boolean functions. The logical expansions have great differences in the complexity of the same function in the three algebra systems. This leads to differences of circuit realization in the circuit area, power consumption, and circuit speed. Detection of the symmetric variables in different types of algebra is very important. In this paper, we propose a new symmetry detection algorithm for identifying 12 types of symmetric variables based on the OR-NXOR expansion. The algorithm has been implemented in C language. Experimental results show that the proposed algorithm is convenient and efficient.

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Table 7 Six methods for comparison

Detection method	Applicability of the number of logical variables	Number of detection types	Foundation work	Time for condition judgment
Graphic method	Usually $n \leq 6$	12	Draw the decomposition charts of $n$ -variable Boolean function	$6n(n-1)$
Spectral coefficients method	Usually $n \leq 6$	12	Calculate the spectral coefficients, then separate the spectral coefficients into $n(n-1)/2$ groups	$6n(n-1)$
Tabular method	Usually $n \leq 10$	2	Draw the 1-valued minterm table	$9n(n-1)/2$
CRM expansion coefficients method	Usually $n \leq 6$	12	Separate the AND-XOR expansion coefficients into $n(n-1)/2$ groups	$6n(n-1)$
$d_j$ -map method	Usually $n \leq 6$	1	Draw the $d_j$ -map	$n(n-1)/2$
Proposed method	No limit	12	Input the number of logical variable in the test function and the OR-NXOR expansion coefficients	$3n(n-1)$

$n$ : the number of logical variables in the test function



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